CHAIR OF APPLIED COMPUTER SCIENCE III

# UNIVERSITY OF MANNHEIM

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Database Systems II Spring Semester 2019

Solution to Exercise Sheet 5

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### Exercise 1

In this exercise, we review some bit manipulation techniques.

Exercise 1 a)

Perform the following bit computations by hand:

- (i) 0110 + 0010
- (ii) 0011 \* 0101
- (iii) 1101 >> 2

Solution

(i) 1110 + 0010 11 -----1000 (ii) 0011 \* 0101 -----0011 + 0011 -----1111

(iii) 1101 >> 2 = 0011, corresponds to an integer division by 4.

## Exercise 1 b)

Explain the two's complement. What is the sum of a positive number and its two's complement?

Solution

In a computer, a negative number is usually represented by the two's complement of its unsigned value. Let's consider the case for -3 (i. e., the representation of negative 3) using a 4-bit integer. The first bit corresponds to the sign, where a 1 is used to represent negative numbers. In order to calculate the negative representation of 3, you need to compute the complement of 3 with respect to  $2^N$ , where N is number of bits, that is, in our example,  $2^4 = 16$ . Therefore, the complement of 3 with respect to 16 is 13 (since 3 + 13 = 16). The binary representation of 13 is 1101. 1101 interpreted as a 4-bit signed integer is therefore -3.

Since the two's complement encodes the negated value of an integer i, we have for every positive integer i that i + twoComplement(i) = i + (-i) = 0.

Consider the example:

	0011	(=	3)
+	1101	(=	-3)
1	10000		

The carry bit is set, but ignored since we only consider N = 4 bits. Therefore, (1)0000 = 0000 (decimal 0), which is the expected result: 3 + (-3) = 0.

Exercise 1 c)

What does the following code do, given n is an integer? ((n & (n-1)) == 0)

#### Solution

This expression evaluates to true if n is a power of 2, otherwise the expression evaluates to false. To see that, note that the expression computes the bitwise and of n and n - 1. If n and n - 1 do not share a 1 bit in any of their positions then ((n & (n-1)) == 0). However, all numbers share a 1 bit in some position with their next smaller value, except powers of 2, since they consist of a single 1-bit and all 0s in the lower significant bits.

Exercise 1 d)

This weeks exercise zip archive contains a file bitvector/bitvector.cc. Implement the setBit and the hasZeroBit member functions of the Bitvector class.

Solution

See code.

Exercise 1 e)

Take a look at the built-in functions that the GCC compiler has to offer. You'll find useful bit manipulation instructions among them.

https://gcc.gnu.org/onlinedocs/gcc/x86-Built-in-Functions.html https://gcc.gnu.org/onlinedocs/gcc/Other-Builtins.html

Exercise 2

Let us consider a database with the following schema.

- Customers: {[<u>id:int</u>, name:char(30), discount:double, country:int]}
- Countries: {[id:int, name:char(30), tax:double]}
- Products: {[<u>id:int</u>, name:char(30), price:double]}
- Orders: {[<u>id:int</u>, customer:int, product:int, quantity:int, date:int, totalPrice:double]}

Exercise 2 a)

Recall the storage layout variant row store and column store from the script.

- (i) Represent the database relations from the above schema in row store layout.
- (ii) Represent the database relations from the above schema in column store layout.

You do not have to write C++ code. Pseudocode that shows the main difference with respect to data organization and data structures is sufficient.

## Solution

Only the solution for relation Countries follows, but the others work similarly, cf. also solution source code.

#### Row store

```
struct country_t {
    int _id;
    char _name[30];
    double _tax;
  }
  std::vector<country_t> countries;
```

## Column store

```
struct Countries {
  std::vector<int> _ids;
  std::vector<std::array<char, 30>> _names;
  std::vector<double> _taxes;
  }
```

```
Exercise 2 b)
```

Download this exercise's zip archive from the website. The folder mmdb contains code that you are asked to complete. The following files are included:

- In common, you find a data generator that creates data with a schema as described above, as well as the basic classes representing customers, countries, products and orders (common/types.hh).
- In rowStore, you find a class RSDatabase that implements a simple row store.
- Additionally, in rowStore, you find the file rsMain.cc that contains a main function and orchestrates the flow of the program for the row store.

Implement a column store for the above schema in a class CSDatabase. You may use the RSDatabase as an orientation. You can use the provided makefile to build the row store database. Warnings like warning: suggest braces around initialization of subobject [-Wmissing-braces] can be ignored.

If you would only like to implement the SQL queries in the next sub-task, the zip archive does also contain an implementation for CSDatabase.

Solution

See code.

Exercise 2 c)

Implement the following the SQL queries for both the row store and the column store. Variables preceded by an \$ represent parameters, i.e. only these parts of the query must be changeable, the rest can be hard-coded. Hint: Implement each query as a member function of the RSDatabase and CSDatabase class.

```
    select totalPrice from
orders
    order by totalPrice desc
fetch first 10 rows only;
```

- select date, sum(totalPrice) from orders where date >= \$date group by date;
- select c.id, c.name, count(o.id) from customers c, orders o where c.id = o.customer group by c.id, c.name;
- update orders
   set totalPrice = \$totalPrice
   where id = \$orderId;

Solution

See code.